

Automatic Generation of HCCA Resistant Scalar Multiplication Algorithm by Proper Sequencing of Field Multiplier Operands

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- Introduction



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- Motivation



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- Public Key Cryptography (PKC) was introduced to address key issues of Key Distribution Problem and Digital Signature Verification problems.
- The two most widely used primitives of PKC are RSA and Elliptic Curve Cryptography.
- Elliptic Curve Cryptography (ECC) has emerged as a strong alternative to RSA due to its property of more security per key bit.



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- Scalar Blinding, Point Coordinate Randomization: Countermeasure against Differential Power Analysis



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- Thus imposes a serious threat to ECC implementations.



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- HCCA threatens an atomic scheme ECC algorithm or unified ECC algorithm (Edward curve) with SPA, DPA resistance.



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- *property 2*: When a pair of multiplications (m_i, m_j) share no common operand among themselves. For example: $(A \times B, C \times D)$
- *property 3*: Given a set S of n field multiplications (m_1, m_2, \dots, m_n) , if there exists at least one pair (m_i, m_j) , where m_i and $m_j \in S$, $i \neq j$, sharing property 1.



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- *condition 1*: Only one of the sets set_a and set_d satisfies property 3.



HCCA scenario 1

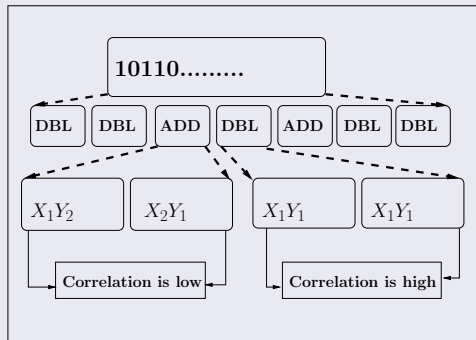


Figure: HCCA scenario 1



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HCCA scenario 2

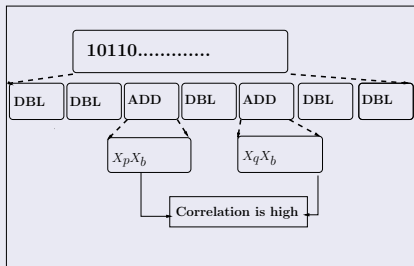


Figure: HCCA scenario 2



Long Integer Multiplication

Algorithm 1: Long Integer Multiplication algorithm(LIM)

Data: : $\{X = (X[t], X[t-1], \dots, X[1])_{2^w}\}$, $\{Y = (Y[t], Y[t-1], \dots, Y[1])_{2^w}\}$

Result: : $\{X.Y\}$

begin

for $i \leftarrow 1$ to $2t$ **do**
 $R[i] = 0$

end

for $i \leftarrow 1$ to t **do**
 $C = 0$;

for $j \leftarrow 1$ to t **do**
 $(U, V)_{2^w} = X[i] \times Y[j]$;
 $(U, V)_{2^w} = (U, V)_{2^w} + C$;
 $(U, V)_{2^w} = (U, V)_{2^w} + R[i+j-1]$;
 $R[i+j-1] = V$;
 $C = U$;

end

$R[i+t] = C$;

end

return R ;

end



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- The leakage $I(O_i)$ is approximated by the Hamming Weight power model.
- A long integer multiplication $LIM(A, B)$ leads to a leakage vector $\langle I_{(a_0b_0)}, I_{(a_0b_1)}, \dots, I_{(a_ib_j)}, \dots, I_{(a_{t-1}b_{t-1})} \rangle$



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 $\text{cov}(\text{LIM}(A, B), \text{cov}(\text{LIM}(C, B))) \neq \text{cov}(\text{LIM}(A, B), \text{LIM}(B, C)).$



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- *Observation 1:* $\rho_1 \neq \rho_2$
- *Observation 2:* $\rho_2 \approx \rho_3$
- *Observation 3:* $\rho_1 > \rho_2$, when $C=A$ (i.e. both the operands are shared).

Countermeasure Design: Safe Sequence

Safe sequence formation for Edward curve formula

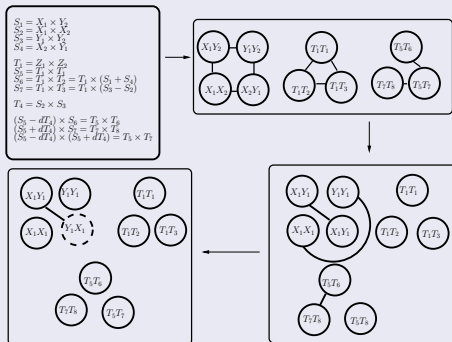


Figure: Safe sequence transformation of Edward unified formula

Safe sequence formation for Brier-Joye unified formula

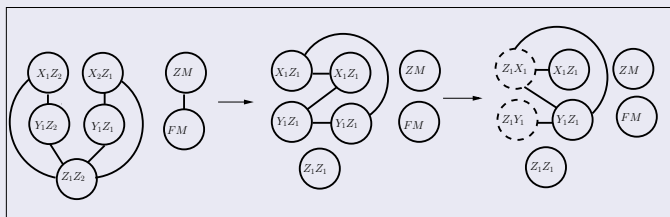


Figure: Safe sequence transformation of Brier-Joye unified formula



- Equipments:



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 - SASEBO GII Board



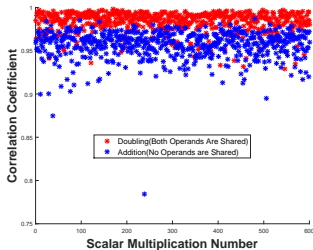
- Equipments:
 - SASEBO GII Board
 - Oscilloscope (DPO4034B)



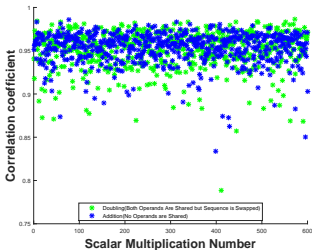
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 - EM Probe



(a) Evaluation of HCCA on Edwards Curve Scalar Multiplier



(b) Evaluation of proposed countermeasure on Edwards Curve Scalar Multiplier



- HCCA scenario 2: Same input point is used in all addition steps
- Re-randomization: Use randomize input point at each stage of addition steps
- After the end of scalar multiplication loop, de-randomize the results¹.
- Similar re-randomization can be used to mitigate other single trace collision attacks ².

¹Poulami Das, Debapriya Basu Roy, Debdeep Mukhopadhyay: Exploiting the Order of Multiplier Operands: A Low Cost Approach for HCCA Resistance. IACR Cryptology ePrint Archive 2015: 925 (2015)

²N. Hanley, H. Kim, and M. Tunstall, Exploiting collisions in addition chain-based exponentiation algorithm using a single trace, Cryptography ePrint Archive: Report 2012/485



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- We show how a unified addition (doubling) formula can be converted into a safe sequence where, the information leakage from sharing of operands among field multipliers have been hidden. Once the sequence have been determined through Algorithm 1 there is no runtime overhead requirement for the step 1 of our countermeasure.
- We have validated HCCA and our proposed countermeasure scheme on a SASEBO platform.

Thank You